

# The Freeness of Ish Arrangements

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First of all, we define three types of arrangements in the  $\ell$ -dimensional Euclidean space  $E$  with coordinates  $x_1, \dots, x_\ell$ :

- Coxeter arrangement  $\text{Cox}(\ell)$
- Shi arrangement  $\text{Shi}(\ell)$
- Ish arrangement  $\text{Ish}(\ell)$

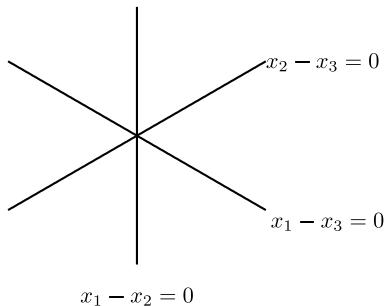
# Three types of arrangement

## Coxeter arrangement

$$\text{Cox}(\ell) = \{ \{x_i - x_j = 0\} \mid 1 \leq i < j \leq \ell \}$$

the **Coxeter arrangement** of the type  $A_{\ell-1}$

Cox(3)



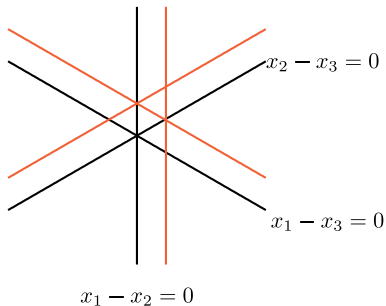
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## Shi arrangement

$$\text{Shi}(\ell) = \text{Cox}(\ell) \cup \{ \{x_i - x_j = 1\} \mid 1 \leq i < j \leq \ell \}$$

the **Shi arrangement** of the type  $A_{\ell-1}$

### Shi(3)



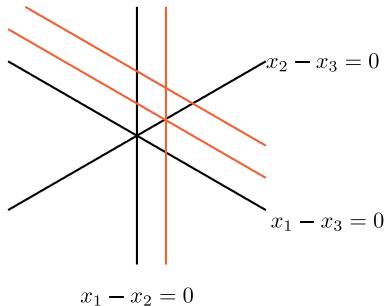
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## Ish arrangement

$$\text{Ish}(\ell) = \text{Cox}(\ell) \cup \left\{ \{x_i - x_j = i\} \mid 1 \leq i < j \leq \ell \right\}$$

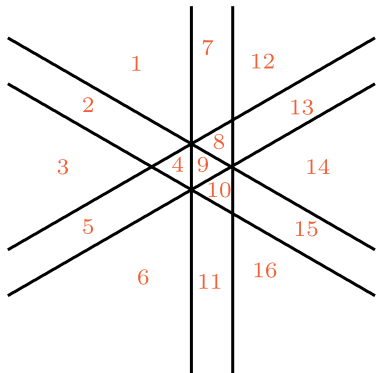
the **Ish arrangement** (of the type  $A_{\ell-1}$ )

### Ish(3)

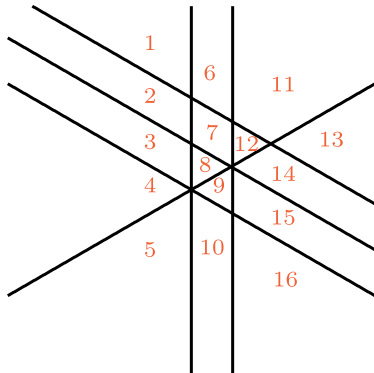


# The number of chambers

$\text{Shi}(\ell)$  and  $\text{Ish}(\ell)$  have the same number of chambers



16 chambers



16 chambers

# The number of chambers

## Definition

Let  $\mathcal{A}$  be an arrangement. Define the **intersection poset**  $L(\mathcal{A})$  by

$$L(\mathcal{A}) = \left\{ \bigcap_{H \in \mathcal{B}} H \neq \emptyset \mid \mathcal{B} \subset \mathcal{A} \right\}, \quad X \leq Y \Leftrightarrow Y \subseteq X,$$

the **Möbius function**  $\mu_{\mathcal{A}} : L(\mathcal{A}) \rightarrow \mathbb{Z}$  by

$$\mu_{\mathcal{A}}(E) = 1, \quad \mu_{\mathcal{A}}(X) = - \sum_{\substack{Y \in L(\mathcal{A}) \\ E \leq Y < X}} \mu_{\mathcal{A}}(Y) \quad (X \neq E).$$

and the **characteristic polynomial**  $\chi(\mathcal{A}, t)$  by

$$\chi(\mathcal{A}, t) = \sum_{X \in L(\mathcal{A})} \mu_{\mathcal{A}}(X) t^{\dim X}.$$

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# The number of chambers

## Theorem (T. Zaslavsky, 1975)

Let  $\mathcal{A}$  be an arrangement in  $E = \mathbb{R}^\ell$ . Then the number of chambers of  $\mathcal{A}$  is  $|\chi(\mathcal{A}, -1)|$ .

## Theorem (D. Armstrong, 2011)

The characteristic polynomial of the Shi arrangement is equal to the characteristic polynomial of the Ish arrangement:

$$\chi(\text{Shi}(\ell), t) = \chi(\text{Ish}(\ell), t) = t(t - \ell)^{\ell-1}.$$

Hence the number of chambers of the Shi arrangement and Ish arrangement are the same.

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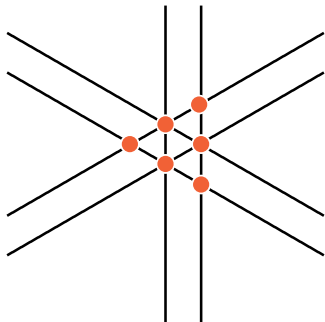
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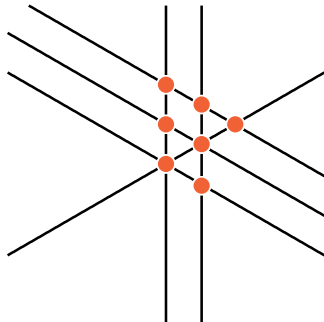
# The intersection posets of Shi and Ish

## Remark

$\text{Shi}(\ell)$  and  $\text{Ish}(\ell)$  share the same characteristic polynomial, although  $L(\text{Shi}(\ell)) \neq L(\text{Ish}(\ell))$ .



6 intersection points



7 intersection points

# Deleted Shi and Ish arrangements

Let  $K_\ell := \{(i, j) \mid 1 \leq i < j \leq \ell\}$  be the edge set of the complete graph with vertices  $\{1, 2, \dots, \ell\}$ .

## Definition

Let  $G \subseteq K_\ell$  be a subgraph. Define the **deleted Shi and Ish arrangements** as follows:

$$\text{Shi}(G) := \text{Cox}(\ell) \cup \left\{ \{x_i - x_j = 1\} \mid (i, j) \in G \right\}$$

$$\text{Ish}(G) := \text{Cox}(\ell) \cup \left\{ \{x_1 - x_j = i\} \mid (i, j) \in G \right\}$$

e.g.

$$\text{Shi}(K_\ell) = \text{Shi}(\ell), \quad \text{Shi}(\emptyset) = \text{Cox}(\ell)$$

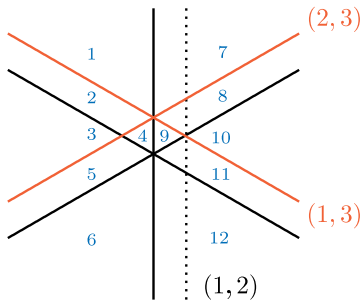
$$\text{Ish}(K_\ell) = \text{Ish}(\ell), \quad \text{Ish}(\emptyset) = \text{Cox}(\ell)$$

# Deleted Shi and Ish arrangements

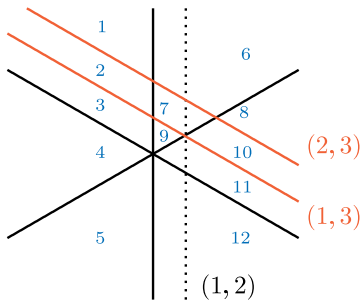
We have the following correspondence between  $\text{Shi}(G)$  and  $\text{Ish}(G)$ :

$$\begin{array}{ccccc}
 \text{Shi}(G) \setminus \text{Cox}(\ell) & \leftrightarrow & G & \leftrightarrow & \text{Ish}(G) \setminus \text{Cox}(\ell) \\
 \Psi & & \Psi & & \Psi \\
 \{x_i - x_j = 1\} & \leftrightarrow & (i, j) & \leftrightarrow & \{x_1 - x_j = i\}
 \end{array}$$

$$G = \{(1, 3), (2, 3)\} \subseteq K_3$$



12 chambers



12 chambers

# Deleted Shi and Ish arrangements

Theorem (D. Armstrong and B. Rhoades, 2012)

*The characteristic polynomial of  $\text{Shi}(G)$  is equal to the characteristic polynomial of  $\text{Ish}(G)$ :*

$$\chi(\text{Shi}(G), t) = \chi(\text{Ish}(G), t).$$

# The freeness of arrangements

From now on, we will consider the problem of whether  $\text{Ish}(G)$  is a free arrangement or not.

## Definition

Let  $S = \mathbb{K}[x_1, \dots, x_\ell]$ . Define the *derivation module*  $\text{Der}_{\mathbb{K}}(S)$  by

$$\text{Der}_{\mathbb{K}}(S) = \{ \theta : S \rightarrow S \mid \theta : \mathbb{K}\text{-linear}, \\ \text{for any } f, g \in S, \theta(fg) = f\theta(g) + \theta(f)g \}.$$

Then  $\text{Der}_{\mathbb{K}}(S)$  is an  $S$ -module.



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# The freeness of arrangements

## Definition

For a hyperplane  $H$ , choose  $\alpha_H \in S$  such that  $\{\alpha_H = 0\} = H$ . For an arrangement  $\mathcal{A}$  of hyperplanes, the polynomial  $Q(\mathcal{A}) = \prod_{H \in \mathcal{A}} \alpha_H \in S$  is called a **defining polynomial**. For a **central** arrangement  $\mathcal{A}$ , (i.e.  $\bigcap_{H \in \mathcal{A}} H \neq \emptyset$ ) define the **logarithmic derivation module**  $D(\mathcal{A})$  of  $\mathcal{A}$  by

$$\begin{aligned} D(\mathcal{A}) &= \{\theta \in \text{Der}_{\mathbb{K}}(S) \mid \theta(Q(\mathcal{A})) \in Q(\mathcal{A})S\} \\ &= \{\theta \in \text{Der}_{\mathbb{K}}(S) \mid \theta(\alpha_H) \in \alpha_H S \text{ for any } H \in \mathcal{A}\}. \end{aligned}$$

We say that  $\mathcal{A}$  is **free**, if  $D(\mathcal{A})$  is a free  $S$ -module.

When  $\mathcal{A}$  is free, there exists a homogeneous basis  $\{\theta_1, \dots, \theta_\ell\}$ , and the multiset  $\text{exp}(\mathcal{A}) = (\deg \theta_1, \dots, \deg \theta_\ell)$  is called the **exponent** of  $\mathcal{A}$ .

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An arrangement  $\mathcal{A}$  is called **supersolvable** if there exists a filtration  $\mathcal{A} = \mathcal{A}_\ell \supseteq \mathcal{A}_{\ell-1} \supseteq \cdots \supseteq \mathcal{A}_1$  such that

- 1  $\text{codim} \bigcap_{H \in \mathcal{A}_i} H = i$  ( $i = 1, 2, \dots, \ell$ )
- 2 For any  $H, H' \in \mathcal{A}_i$  with  $H \neq H'$ , there exists some  $H'' \in \mathcal{A}_{i-1}$  such that  $H \cap H' \subseteq H''$ .

Cox( $\ell$ ) is supersolvable

Cox( $k$ ) =  $\{ \{x_i - x_j = 0\} \mid 1 \leq i < j \leq k \}$  ( $2 \leq k \leq \ell$ ) can be regarded as arrangements in  $\mathbb{R}^{\ell-1}$ . The filtration

$$\text{Cox}(\ell) \supseteq \text{Cox}(\ell - 1) \supseteq \cdots \supseteq \text{Cox}(2)$$

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# The freeness of arrangements

## Theorem (Addition theorem, H. Terao, 1980)

Let  $\mathcal{A}$  be a central arrangement and  $H_0 \in \mathcal{A}$ . If  $\mathcal{A}' := \mathcal{A} \setminus \{H_0\}$  and  $\mathcal{A}'' := \{H_0 \cap H \mid H \in \mathcal{A}'\}$  are free and  $\exp \mathcal{A}'' \subset \exp \mathcal{A}'$ , then  $\mathcal{A}$  is free.

## Definition

Define the *inductive freeness* by the following:

- 1 The empty arrangement is inductively free.
- 2  $\mathcal{A}$  is inductively free if there exists  $H_0 \in \mathcal{A}$  such that  $\mathcal{A}'$  and  $\mathcal{A}''$  are inductively free and  $\exp \mathcal{A}'' \subset \exp \mathcal{A}'$ .

## Theorem

*supersolvable  $\Rightarrow$  inductively free  $\Rightarrow$  free*



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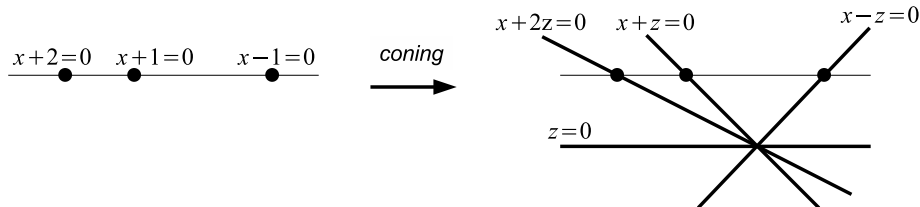
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## Definition

Let  $\mathcal{A}$  be an *affine* arrangement in  $E$  and  $E^* = \langle x_1, \dots, x_\ell \rangle_{\mathbb{R}}$ . Let  $W$  be a vector space such that  $W^* = \langle x_1, \dots, x_\ell, z \rangle_{\mathbb{R}}$ . Then we define a *central* arrangement  $\mathbf{c}\mathcal{A}$  in  $W$  by

$$Q(\mathbf{c}\mathcal{A}) = z \cdot z^{\deg Q(\mathcal{A})} Q(\mathcal{A}) \left( \frac{x_1}{z}, \dots, \frac{x_\ell}{z} \right).$$

$\mathbf{c}\mathcal{A}$  is called the *cone* over  $\mathcal{A}$ .



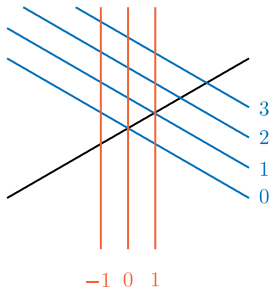
# $N$ -Ish arrangement

## Definition

Let  $N = (N_2, N_3, \dots, N_\ell)$  with finite sets  $N_j \subseteq \mathbb{R}$  for  $2 \leq j \leq \ell$ . Define the  **$N$ -Ish arrangement**  $\text{Ish}(N)$  by

$$\text{Ish}(N) = \left\{ \{x_i - x_j = 0\} \mid 2 \leq i < j \leq \ell \right\} \cup \left\{ \{x_1 - x_j = a\} \mid 2 \leq j \leq \ell, a \in N_j \right\}.$$

$\text{Ish}(N) \quad N = (\{-1, 0, 1\}, \{0, 1, 2, 3\})$



# $N$ -Ish arrangement

The class of  $N$ -Ish arrangements includes the deleted Ish arrangements.

$$N_j = \{0, 1, \dots, j-1\} \Rightarrow \text{Ish}(N) = \text{Ish}(\ell)$$

$$N_j = \{0\} \cup \{i \mid (i, j) \in G\} \Rightarrow \text{Ish}(N) = \text{Ish}(G)$$

## Definition

$N = (N_2, \dots, N_\ell)$  is a **nest** if there exists a permutation  $w$  of  $\{2, 3, \dots, \ell\}$  such that  $N_{w(2)} \subseteq N_{w(3)} \subseteq \dots \subseteq N_{w(\ell)}$

$N_j = \{0, 1, \dots, j-1\} \Rightarrow N_2 \subseteq N_3 \subseteq \dots \subseteq N_\ell \Rightarrow N = (N_2, \dots, N_\ell)$  is a nest

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# *N*-Ish arrangement

## Theorem (Abe, S, Tsujie)

*For  $N$ -Ish arrangements, the following are equivalent:*

- ①  *$N$  is a nest,*
- ②  *$\mathbf{cIsh}(N)$  is supersolvable,*
- ③  *$\mathbf{cIsh}(N)$  is inductively free,*
- ④  *$\mathbf{cIsh}(N)$  is free.*

(①  $\Rightarrow$  ②) Let  $N_2 \supseteq N_3 \supseteq \cdots \supseteq N_\ell$ ,

$$\mathcal{I}_i := \{ \{x_1 - x_j = az\} \mid 2 \leq j \leq i, a \in N_j \}$$

$$\cup \{ \{x_j - x_k = 0\} \mid 2 \leq j < k \leq i \} \cup \{ \{z = 0\} \}.$$

Then  $\mathbf{cIsh}(N) = \mathcal{I}_\ell \supseteq \mathcal{I}_{\ell-1} \supseteq \cdots \supseteq \mathcal{I}_1$ .

## Corollary

*$\mathbf{cIsh}(\ell)$  is a free arrangement.*



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- ④  *$\mathbf{cIsh}(N)$  is free.*

(①  $\Rightarrow$  ②) Let  $N_2 \supseteq N_3 \supseteq \cdots \supseteq N_\ell$ ,

$$\mathcal{I}_i := \{ \{x_1 - x_j = az\} \mid 2 \leq j \leq i, a \in N_j \}$$

$$\cup \{ \{x_j - x_k = 0\} \mid 2 \leq j < k \leq i \} \cup \{ \{z = 0\} \}.$$

Then  $\mathbf{cIsh}(N) = \mathcal{I}_\ell \supseteq \mathcal{I}_{\ell-1} \supseteq \cdots \supseteq \mathcal{I}_1$ .

## Corollary

*$\mathbf{cIsh}(\ell)$  is a free arrangement.*

# *N*-Ish arrangement

## Theorem (Abe, S, Tsujie)

For *N*-Ish arrangements, the following are equivalent:

- ① *N* is a nest,
- ②  $\mathbf{cIsh}(N)$  is supersolvable,
- ③  $\mathbf{cIsh}(N)$  is inductively free,
- ④  $\mathbf{cIsh}(N)$  is free.

(①  $\Rightarrow$  ②) Let  $N_2 \supseteq N_3 \supseteq \cdots \supseteq N_\ell$ ,

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## Corollary

$\mathbf{cIsh}(\ell)$  is a free arrangement.

# *N*-Ish arrangement

## Theorem (Abe, S, Tsujie)

Suppose that  $N_2 \subseteq \cdots \subseteq N_\ell$ . Define homogeneous derivations  $\theta_0, \theta_1, \dots, \theta_\ell$  by

$$\theta_0 := \sum_{i=1}^{\ell} \partial x_i, \quad \theta_1 = \sum_{i=1}^{\ell} x_i \frac{\partial}{\partial x_i} + z \frac{\partial}{\partial z},$$
$$\theta_k := \sum_{s=2}^k \left( \prod_{a \in N_k} (x_1 - x_s - az) \prod_{t=k+1}^{\ell} (x_s - x_t) \right) \frac{\partial}{\partial x_s} \quad (2 \leq k \leq \ell).$$

Then  $\theta_0, \theta_1, \dots, \theta_\ell$  form a basis for  $D(\mathbf{cIsh}(N))$ . In particular, the exponents are given by

$$\exp(\mathbf{cIsh}(N)) = (0, 1, |N_2| + \ell - 2, |N_3| + \ell - 3, \dots, |N_\ell|).$$

# The characteristic polynomial of the Ish arrangement

## Theorem (Factorization Theorem, H. Terao, 1981)

Let  $\mathcal{A}$  be a free arrangement with  $\exp(\mathcal{A}) = \{d_1, \dots, d_\ell\}$ . Then the characteristic polynomial of  $\mathcal{A}$  factors as  $\chi(\mathcal{A}, t) = \prod_{i=1}^{\ell} (t - d_i)$ .

The characteristic polynomial  $\chi(\text{Ish}(\ell), t)$

$\text{Ish}(\ell)$  is  $\text{Ish}(N)$  with  $N_j = \{0, 1, \dots, j - 1\}$ .

Since  $|N_j| = j$ ,

$$\begin{aligned}\exp(\mathbf{cIsh}(\ell)) &= (0, 1, |N_2| + \ell - 2, |N_3| + \ell - 3, \dots, |N_\ell|) \\ &= (0, 1, \ell, \dots, \ell).\end{aligned}$$

Hence  $\chi(\mathbf{cIsh}(\ell), t) = t(t - 1)(t - \ell)^{\ell - 1}$ . Since  $\chi(\mathbf{c}\mathcal{A}, t) = (t - 1)\chi(\mathcal{A}, t)$ ,

$$\chi(\text{Ish}(\ell), t) = t(t - \ell)^{\ell - 1}.$$

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## Theorem (C. A. Athanasiadis, 1998)

Let  $G \subseteq K_\ell$ . Then Shi( $G$ ) is free if and only if the following condition (A) is satisfied:

- (A) *There exists a permutation  $w$  of  $\{1, 2, \dots, \ell\}$  such that  $w^{-1}G \subseteq K_\ell$  and*
- $$1 \leq i < j < k \leq \ell, (i, j) \in w^{-1}G \Rightarrow (i, k) \in w^{-1}G.$$

# Shi( $G$ ) and Ish( $G$ )

For  $G \subseteq K_\ell$ , define  $N_G = (N_2, \dots, N_\ell)$  by

$$N_j = \{0\} \cup \{i \mid (i, j) \in G\}.$$

Then  $\text{Ish}(G) = \text{Ish}(N_G)$ .

## Lemma

*$G$  satisfies the condition (A)  $\Leftrightarrow N_G$  is a nest*

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## Lemma

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## Theorem (Abe, S, Tsujie)

Let  $G \subseteq K_\ell$ . Then the following are equivalent.

- 1  $G$  satisfies (A)
- 2  $N_G$  is a nest
- 3  $\text{cShi}(G)$  is free
- 4  $\text{cIsh}(G)$  is free

# Problem

## Problem

*Is there an isomorphism between  $D(\mathbf{cShi}(G))$  and  $D(\mathbf{cIsh}(G))$  when they are not free?*

Thank you very much  
for your attention!